

MODELS OF SET THEORY

IN GENERAL A SET (OR CLASS) Π WITH A BINARY RELATION E CAN SERVE AS A STRUCTURE FOR THE LANGUAGE OF SET THEORY

WE INTERPRET FORMULAS AS FOLLOWS

$$(x \in y)^{\Pi, E} \text{ IS } x \in y$$

$$(x = y)^{\Pi, E} \text{ IS } x = y$$

$$(\neg \varphi)^{\Pi, E} \text{ IS } \neg \varphi^{\Pi, E}$$

$$(\varphi \wedge \psi)^{\Pi, E} \text{ IS } \varphi^{\Pi, E} \wedge \psi^{\Pi, E}$$

$$(\exists x)(\varphi)^{\Pi, E} \text{ IS } (\exists x \in \Pi)(\varphi)^{\Pi, E}$$

(THIS SUFFICES: \rightarrow , \vee , \Leftrightarrow AND \forall CAN BE EXPRESSED IN TERMS OF \neg , \wedge AND \exists)

USE: TO SHOW THAT CERTAIN FORMULAS ARE NOT FORMALLY DERIVABLE FROM OTHERS

LANGUAGE OF GROUPS: $*$, \cdot^{-1} , e

INTERPRET $*$ AS $+$; \cdot^{-1} AS $-$; e AS 0
IN \mathbb{Z}

HERE $(\forall x)(\forall y)(x * y = y * x)$ IS TRUE

INTERPRET $*$ AS \circ ; \cdot^{-1} AS \cdot^{-1} ; e AS (1)
IN S_2

HERE $(\forall x)(\forall y)(x * y = y * x)$ IS FALSE

NEITHER THAT FORMULA NOR ITS NEGATION IS DERIVABLE FROM THE GROUP AXIOMS

EXAMPLE $\mathcal{M} = \mathbb{Z}$; E IS $<$.

$((\forall x)(\exists y)(x < y))^{>, <}$ THEN IS

$$(\forall x \in \mathbb{Z})(\exists y \in \mathbb{Z})(x < y)$$

$$((\forall x)(\forall y)(\exists z)(\forall u)(u < z \leftrightarrow (u = x \vee u = y)))^{>, <}$$

BECOMES

$$(\forall x \in \mathbb{Z})(\forall y \in \mathbb{Z})(\exists z \in \mathbb{Z})(\forall u \in \mathbb{Z})(u < z \leftrightarrow (u = x \vee u = y))$$

NOT TRUE SO AXIOM OF PAIRING DOES NOT HOLD IN THIS STRUCTURE

AXIOM OF UNION:

$$(\forall x)(\exists y)(\forall z)(z \in y \leftrightarrow (\exists u \in x)(z \in u))$$

BECOMES

$$(\forall x \in \mathbb{Z})(\exists y \in \mathbb{Z})(\forall z \in \mathbb{Z})(z < y \leftrightarrow (\exists u \in x)(z < u \vee z = u))$$

NOTE $(\exists u \in x)(\dots)$

$$\text{IS } (\exists u \in x)(u \in x \wedge \dots)$$

THIS IS TRUE: TAKE $y = x - 1$

SO AXIOM OF UNION DOES NOT IMPLY THE AXIOM OF PAIRING.

HOMEWORK CHECK POWER SET AND REGULARITY IN THIS STRUCTURE

USUALLY WE LET E BE \in ITSELF

QUITE OFTEN, BUT NOT ALWAYS / $\varphi^{\mathcal{M}, \mathcal{G}}$ IS WRITTEN

\mathcal{M} IS TRANSITIVE.

$\varphi^{\mathcal{M}}$

TO CHECK / VERIFY AXIOMS IN SUCH SITUATIONS IT HELPS IF THINGS DON'T CHANGE WHEN GOING INTO \mathcal{M} .

THINGS LIKE $x = \{y, z\}$, $x = \langle y, z \rangle$, $x = \cup y$,
 $x = \emptyset$, $x \subseteq y$, x IS TRANSITIVE,
 x IS AN ORDINAL, $x = \omega$, ...

ARE ABSOLUTE FOR M :

WE SAY φ IS ABSOLUTE FOR M
 IF φ AND φ^M ARE EQUIVALENT.

Δ_0 -FORMULAS ARE ABSOLUTE FOR
 TRANSITIVE CLASSES

Δ_0 -FORMULAS:

- $x \in y$
- $x = y$
- IF φ IS Δ_0 THEN SO IS $\neg \varphi$.
- IF φ AND ψ ARE THEN SO IS $\varphi \wedge \psi$.
- IF φ IS THEN SO IS $(\exists x \in y)(\varphi)$
 (AND $(\forall x \in y)(\varphi)$)

PROOF: INDUCTION ON COMPLEXITY

- $(x \in y)^M$ IS $(x \in y)$
 - $(x = y)^M$ IS $(x = y)$
 - IF $\varphi \leftrightarrow \varphi^M$ THEN $\neg \varphi \leftrightarrow \neg \varphi^M$
 - IF $\varphi \leftrightarrow \varphi^M$ AND $\psi \leftrightarrow \psi^M$ THEN $\varphi \wedge \psi \leftrightarrow (\varphi \wedge \psi)^M$
- } SIMPLE LOGIC.
- ASSUME $\varphi \leftrightarrow \varphi^M$
- CERTAINLY IF $(\exists x \in M)(\cup x \wedge \varphi^M)$
 THEN $(\exists x \in M)(\cup x \wedge \varphi)$
 AND $(\exists x \in M)(\cup x \wedge \varphi)$ BY IND. ASS.
- CONVERSELY IF $x \in M$ AND $(\cup x \wedge \varphi)$
 THEN, AS $x \in M$, $(\cup x \wedge \varphi^M)$ BY IND. ASS.
 NOW USE $\varphi \leftrightarrow \varphi^M$ TO CONCLUDE
 $(\cup x \wedge \varphi) \leftrightarrow (\cup x \wedge \varphi^M)$

EXAMPLE, BY THE AXIOM OF REGULARITY,

\aleph_α IS AN ORDINAL $\Leftrightarrow \aleph_\alpha$ IS TRANSITIVE
AND LINEARLY ORDERED
BY \in .

THE LATTER IS Δ_0 :

$(\forall y \in x)(\forall z \in x)(y = z \vee y \in z \vee z \in y) \wedge$

$(\forall y \in x)(\forall z \in y)(z \in x)$

Δ_0 -FORMULAS ARE USEFUL BECAUSE:

A TRANSITIVE M SATISFIES

PAIRING IFF $\forall x \in M \forall y \in M \exists z \in M$

UNION IFF $\forall x \in M \cup x \in M$

INFINITY IFF $\omega \in M$

LET $M = \bigcup_{\alpha \in \text{ORD}} V_\alpha$ THEN

ZF^- IMPLIES σ^M FOR ALL AXIOMS σ
OF ZF

\uparrow
ZF MINUS REGULARITY

SO IF ZF^- IS CONSISTENT THEN SO IS ZF .

V_ω SATISFIES ALL AXIOMS OF ZFC
EXCEPT INFINITY

$V_{\omega+\omega}$ SATISFIES ALL AXIOMS OF $ZFC(C)$
EXCEPT REPLACEMENT:

THE WELL-ORDER \triangleleft ON ω
DEFINED BY

$n \triangleleft m$ IF n IS EVEN AND m IS ODD
OR $(m \equiv n \pmod{2})$
AND $n \in m$

~~HAS~~ IS NOT ISOMORPHIC TO ANY
ORDINAL IN $V_{\omega+\omega}$.

IF κ IS INACCESSIBLE

THEN V_κ SATISFIES ALL OF ZFC

IF κ IS THE FIRST INACCESSIBLE

THEN $V_\kappa \models$ "THERE ARE NO INACCESSIBLES"

SO ZFC DOES NOT PROVE INACCESSIBLES EXIST

THERE IS AN INACCESSIBLE

ALSO WE CANNOT PROVE

(*) IF ZFC IS CONSISTENT THEN SO IS ZFC + I

WE JUST SAW

ZFC + I \vdash "ZFC IS CONSISTENT"

MODUS PONENS FROM (A) WOULD GIVE

ZFC + I \vdash "ZFC + I IS CONSISTENT"

CONTRADICTING GÖDEL'S SECOND INCOMPLETENESS THEOREM

FOR A REGULAR κ WE LET

$$H(\kappa) = \{x : |TC(x)| < \kappa\}$$

(THIS SET EXISTS $\because H(\kappa) \in V_\kappa$)

$$H(\omega) = V_\omega$$

$H(\omega_1)$ IS THE SET OF HEREDITARILY COUNTABLE SETS

$H(\kappa)$ SATISFIES ALL OF ZFC EXCEPT POSSIBLY POWER SET

CASE IN POINT:

$H(\omega_1) \models$ "ALL SETS ARE COUNTABLE"

ELEMENTARY SUBSTRUCTURES

THEOREM 12.1 LÖWENHEIM-SKOLEM

EVERY STRUCTURE FOR A COUNTABLE LANGUAGE HAS A COUNTABLE ELEMENTARY SUBSTRUCTURE

\mathcal{M} IS ELEMENTARY SUBSTRUCTURE OF \mathcal{K}
 $\mathcal{M} \prec \mathcal{K}$ MEANS:

FOR EVERY FORMULA φ AND ALL $a_1, \dots, a_n \in \mathcal{M}$:

$$\varphi^{\mathcal{M}}(a_1, \dots, a_n) \Leftrightarrow \varphi^{\mathcal{K}(x)}(a_1, \dots, a_n)$$

EXTENSIONS

- IF $X \subseteq \mathcal{K}$ THEN THERE IS $\mathcal{M} \prec \mathcal{K}$ WITH $|\mathcal{M}| \leq |X| \cdot \aleph_0$
- ONE CAN ALSO ACHIEVE $\aleph_w \subseteq \mathcal{M}$ TOGETHER WITH $|\mathcal{M}| \leq |X| \cdot 2^{\aleph_0}$

LET $\kappa > \omega$, AND LET $\mathcal{M} \prec \mathcal{K}$ BE A COUNTABLE ELEMENTARY SUBSTRUCTURE

NOTE \mathcal{M} IS DEFINITELY NOT TRANSITIVE, (AS WE SHALL SEE MOMENTARILY.)

EASY OBSERVATIONS

LET $x \in \mathcal{M}$ BE NONEMPTY SO

$$(\exists y)(y \in x) \text{ IS TRUE IN } \mathcal{K}$$

SO IT IS TRUE IN \mathcal{M} , THAT IS

$$(\exists y \in \mathcal{M})(y \in x)$$

SO IF $x \in \Pi$ AND $x \neq \emptyset$ THEN $x \cap \Pi \neq \emptyset$.

LIKEWISE: IF $x, y \in \Pi$ AND $x \neq y$
THEN $x \cap \Pi \neq y \cap \Pi$

$(\exists x)(\forall y \in x)(y \neq y)$ HOLDS IN $H(K)$ ($x = \emptyset$)

SO $(\exists x \in \Pi)(\forall y \in \Pi)(y \in x \rightarrow y \neq y)$ IS TRUE

TAKE SUCH AN $x \in \Pi$ THEN $x \cap \Pi = \emptyset$
AND SO $x = \emptyset$.

SO $\emptyset \in \Pi$ (NECESSARILY)

$(\exists x)(\forall y)(y \in x \leftrightarrow y = \emptyset)$ HOLDS IN $H(K)$

SO $(\exists x \in \Pi)(\forall y \in \Pi)(y \in x \leftrightarrow y = \emptyset)$

SO $1 = \{\emptyset\} \in \Pi$ (NECESSARILY)

INDUCTION: $(\forall n \in \omega)(n \in \Pi)$ (EXERCISE)

ALSO $\omega \in \Pi$ USE THE FOLLOWING FORMULA

$(\exists x)(\forall y)(y \in x \rightarrow (y = \emptyset \vee (\exists z \in x)(y = z \cup \{z\})))$

(SEE HOMEWORK PROBLEM 1.8)

THERE IS SUCH AN x IN Π AND IT MUST
BE ω ITSELF

ALSO $\omega_1 \in \Pi$ AS THE FIRST ORDINAL
SUCH THAT THERE IS NO BIJECTION
BETWEEN IT AND ω .

BUT OF COURSE $\omega_1 \notin \Pi$.

LEMMA IF $x \in \Pi$ AND x IS COUNTABLE
THEN $x \subseteq \Pi$.

PROOF $(\exists f)(f: \omega \xrightarrow{\text{ONTO}} x)$ IS TRUE IN $H(K)$

SO $(\exists f \in \Pi)(f: \omega \xrightarrow{\text{ONTO}} x)$ TAKE SUCH AN f .

THEN $(\forall n \in \omega)(\exists y)(\langle n, y \rangle \in f)$ (IN $H(K)$, HENCE IN Π)

SO $(\forall n \in \omega)(\exists y \in \Pi)(\langle n, y \rangle \in f)$

BUT y IS UNIQUE SO $f \cap \Pi \in \Pi$.

WE SEE $\{f \cap \Pi : n \in \omega\} \subseteq \Pi$.

Let $f: \omega_1 \rightarrow \mathbb{R}$ be continuous

Let $\kappa \geq (2^{\aleph_1})^+$ and let $M \subset M(\kappa)$
be countable with $\underline{f} \in M$

By the lemma $\delta_M = \Pi \omega_1$ is an ordinal.

EXERCISE: prove δ_M is an ε -number

We show f is constant on $[\delta_M, \omega_1)$

Let $n \in \omega$ and take $\beta < \delta_M$ such that

$$(\forall \gamma \in (\beta, \delta_M]) (|f(\gamma) - f(\delta_M)| < 2^{-(n+1)})$$

Then $(\forall \gamma, \eta \in (\beta, \delta_M)) (|f(\gamma) - f(\eta)| < 2^{-n})$

by the triangle inequality.

So $(\forall \gamma \in M) (\forall \eta \in M) (\gamma, \eta > \beta \rightarrow |f(\gamma) - f(\eta)| < 2^{-n})$

that is $(\forall \gamma) (\forall \eta) (\gamma, \eta > \beta \rightarrow |f(\gamma) - f(\eta)| < 2^{-n})^M$

is true, and so, because $\beta, \delta, n \in M$

$$((\forall \gamma) (\forall \eta) (\gamma, \eta > \beta \rightarrow |f(\gamma) - f(\eta)| < 2^{-n}))^{M(\kappa)}$$

is true.

This shows: if $\delta, \eta \geq \delta_M$ then

$$|f(\delta) - f(\eta)| < 2^{-n} \text{ for all } n.$$

That is f is constant on $[\delta_M, \omega_1)$

EXERCISE prove that there is a $\delta < \delta_M$
such that f is constant on
 $[\delta, \omega_1)$.

EXERCISE prove: if C is club and $C \in M$
then $\delta_M \in C$.

EXERCISE prove: if $S \in M$ and $\delta_M \in S$
then S is stationary.